

The Hurwitz action on real reflection groups

(w/ Vic Reiner)

G a group

$(t_1, \dots, t_i, t_{i+1}, \dots, t_m) \in G^m$

$\downarrow \sigma_i$

$(t_1, \dots, t_{i+1}, t_{i+1}^{-1} t_i t_{i+1}, \dots, t_m)$

Hurwitz
move

Hurwitz
action

Braid group
action

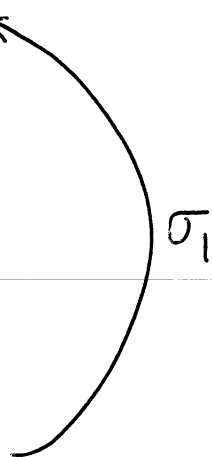
Obs: preserves product $t_1 \dots t_m =: c$, so acts
on factorizations

Question: connectedness? (Equiv: orbit structure.)

(Background:

- Hurwitz, 1890s, $G = \underline{\underline{S_n}}$, t_i transpositions,
something to do w/ covering spaces of $\mathbb{R}e$
Riemann sphere
- Bessis, 2000s, $G = \underline{\text{finite real refln}}$, t_i reflections,
dual Coxeter systems, eventually something to
do w/ $K(\pi, 1)$ Eilenberg-MacLane for complement
of \mathbb{C} hyperplane arrangement of \mathbb{C} refln gps.

E.g. S_3

$$\begin{aligned} C = (123) &= \overset{s_1}{(12)} \overset{s_2}{(23)} \\ &\downarrow \sigma_1 \\ &= (23)(13) \\ &\downarrow \sigma_1 \\ &= (13)(12) \end{aligned}$$


Thm (Bessis; "dual Matsumoto-Tits lemma"): W a finite real refln gp of rank n , $c \in W$ a Coxeter element. The Hurwitz action on minimal (length n) factorizations of c into reflections is transitive.

Thm (L-Reiner) W, c , ditto. $m \in \mathbb{N}$. The Hurwitz action on factorizations of c as a product of m reflections is as transitive as possible.

Finite real refln gp: matrix gp (finite, over \mathbb{R}) generated by its subset of orthogonal reflections thro a hyperplane.

Generators & relations:
 $\langle s_1, \dots, s_n \mid s_i^2 = 1, (s_i s_j)^{m_{ij}} = 1 \rangle$
 ↑
 "simple refls"

Coxeter el't: $c = s_1 s_2 \dots s_n$ (& conjugates)

E.g.: • S_n : perm matrices, $s_i =$ adjacent transpositions,
 $c = (1234 \dots n)$

• type B: hyperoctahedral gp / signed perm. matrices,
 $c = \begin{pmatrix} 1 & 2 & 3 & \dots & n-1 & n \\ 2 & 3 & 4 & \dots & n & -1 \end{pmatrix}$
 ↑ (two-line)
 $= \begin{pmatrix} 1 & 2 & 3 & \dots & n-1 & n \\ 1 & 0 & 0 & \dots & 0 & 0 \\ 0 & 1 & 0 & \dots & 0 & 0 \\ 0 & 0 & 1 & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & 1 & 0 \end{pmatrix}$ ← (as a matrix)

E.g. S_3

$$\begin{aligned} C = (123) &= \overset{s_1}{(12)} \overset{s_2}{(23)} \\ &\downarrow \sigma_1 \\ &= (23)(13) \\ &\downarrow \sigma_1 \\ &= (13)(12) \end{aligned}$$

σ_1

E.g. Type B_2

$$W = \left\{ \begin{bmatrix} \pm 1 & 0 \\ 0 & \pm 1 \end{bmatrix}, \begin{bmatrix} 0 & \pm 1 \\ \pm 1 & 0 \end{bmatrix} \right\} \quad (8 \text{ elements})$$

Reflections

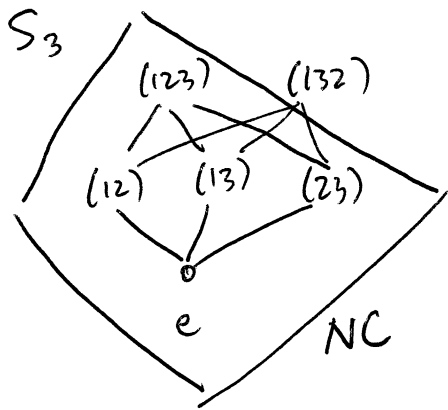
$$\begin{array}{cc} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} & \begin{bmatrix} -1 & 0 \\ 0 & 1 \end{bmatrix} \\ \begin{bmatrix} 0 & -1 \\ -1 & 0 \end{bmatrix} & \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \end{array}$$

$$\begin{aligned} C &= \overset{s_1 s_2}{\begin{bmatrix} 0 & -1 \\ +1 & 0 \end{bmatrix}} = s_1 \cdot s_1 \cdot s_1 \cdot s_2 \\ &= s_2 \cdot s_2 \cdot s_1 \cdot s_2 \end{aligned}$$

An idea from Bessis' pf
 (& ~~related~~ related work):
 view (t_1, \dots, t_m) as a walk

$e \rightarrow t_1 \rightarrow t_1 t_2 \rightarrow \dots \rightarrow t_1 t_2 \dots t_m$
 on W .

Elements we meet on shortest
 paths to c are
 "W-noncrossing partitions"



But if
 $m > n$, we
 don't need to
 stay in
 $NC(W)$!

So we cannot use
 existing tools for NC.

Structure of pf:

Main Thm

↑

Lemma 1: factorizations of
arbitrary $w \in W$

↑

Lemma 2: root circuits are
 acutely disconnected

Lemma 1 (L-Perin) W a finite real reflection group,
 w an arbitrary element of W ,
 $t = (t_1, \dots, t_m)$ a factorization of w into reflections.

The Hurwitz orbit of t contains a factorization

$(t'_1, t'_2, t'_3, t'_4, \dots, t'_{2k-1}, t'_{2k}, t'_{2k+1}, \dots, t'_m)$ s.t.

$t'_1 = t'_2, t'_3 = t'_4, \dots, t'_{2k-1} = t'_{2k}$ & (t'_{2k+1}, \dots, t'_m) is

a shortest factorization of w .

Use this to put long factorizations into semi-canonical form;
then use Bessis' results about shortest factorizations of
Cox. elts to finish.

Root systems

W a finite real refl'n gp.
Take a pair $\pm\alpha$ of vectors
 \perp to the reflecting hyperplane
of each reflection in W .
This is a root system

$$\Phi = \Phi(W).$$

(Also axioms ---)

Lemma 2 (L-Reiner)

Suppose $\{\alpha_1, \dots, \alpha_k\}$ is a circuit
(i.e., minimal dependent set) in a
root system. Replace α_i by $-\alpha_i$ as
needed s.t. dependence

$$c_1\alpha_1 + c_2\alpha_2 + \dots + c_k\alpha_k = 0$$

has $c_i > 0$. Define acuteness graph

$$G \text{ by } V(G) = \{1, \dots, k\},$$

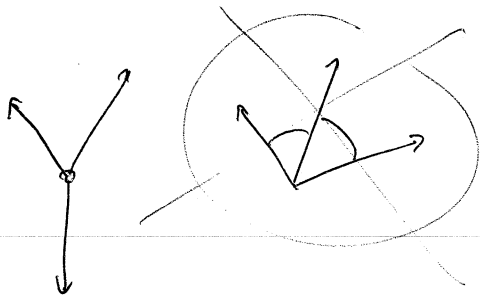
$$E(G) = \{ \{i, j\} : \alpha_i \cdot \alpha_j > 0 \}.$$

Then G is disconnected.

Runk (Fiedler): if α_i not restricted to a root system, obtuseness graph must be connected, acuteness graph is arbitrary.

How to prove acuteness graphs are disconnected?

RK 2:



Classical types A, B, C, D's

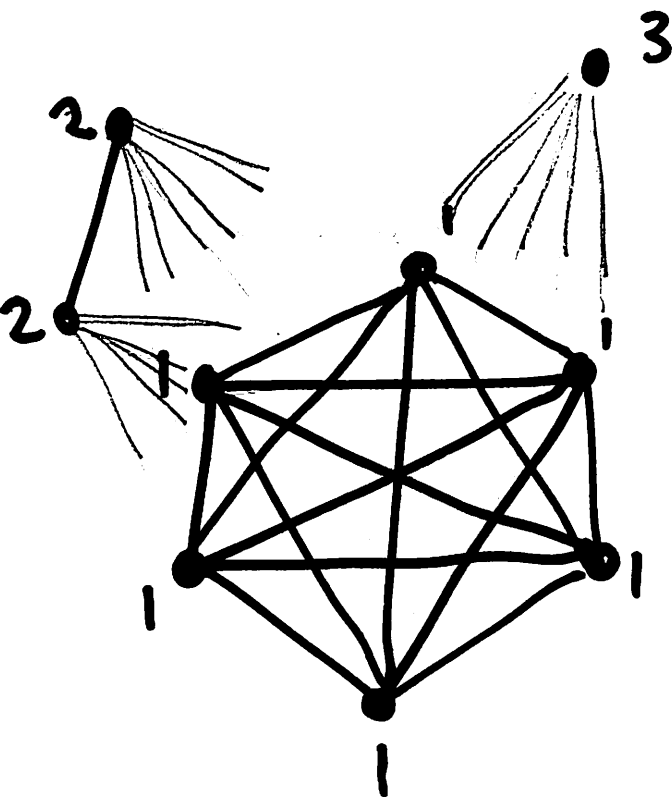
Zaslavsky, signed graphs, circuits are understood, very few edges in acuteness graph.

Exceptional types: weeks of computer computations ".

Lots of nice data produced; interested?

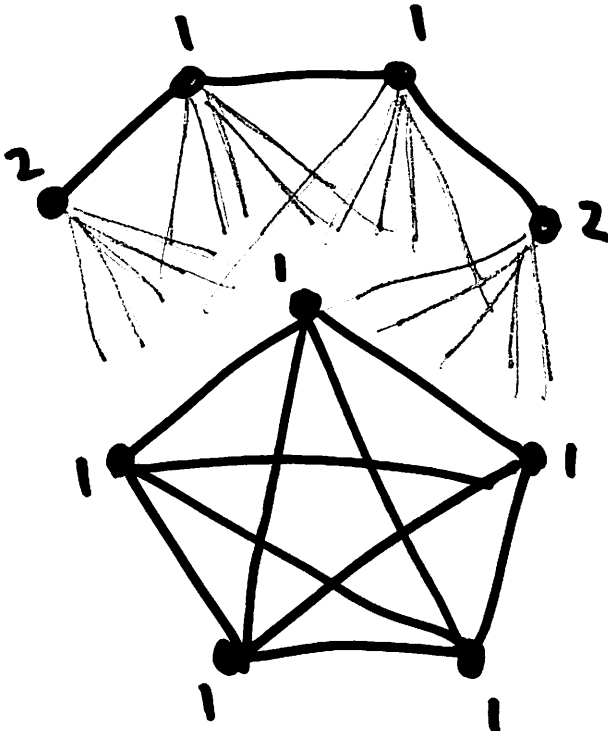
Eg examples

16 acute edges,
3 components



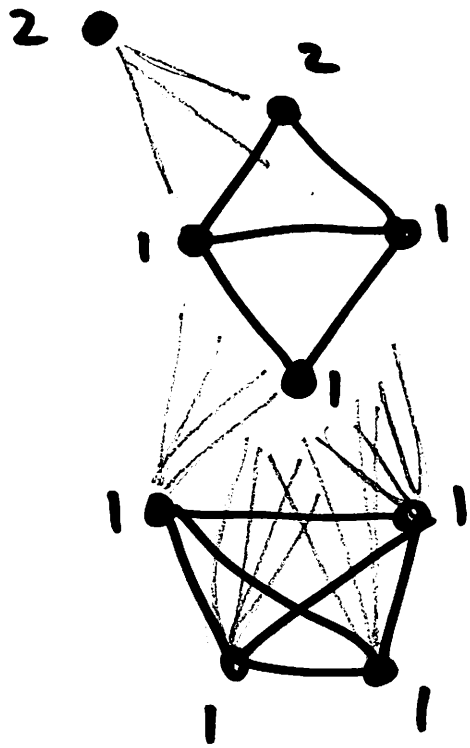
(18 obtuse edges)

13 acute edges,
2 components



(20 obtuse edges)

11 acute edges,
3 components



(19 obtuse edges)